

# On Linial's Conjecture for Split Digraphs

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**Abstract.** *In this paper we show that Linial's Conjecture holds for two classes of split digraphs, namely the spider digraphs and the  $k$ -loose digraphs.*

## 1. Introduction

The digraphs considered in this text do not contain loops or parallel arcs and by path we mean directed path. Let  $D$  be a digraph. We denote by  $V(P)$  the set of vertices of a path  $P$ . The **size** of a path  $P$ , denoted by  $|P|$ , is  $|V(P)|$ <sup>1</sup>. We denote by  $\lambda(D)$  the size of the longest path in  $D$  and by  $\alpha(D)$  the size of a maximum stable set. A *path partition*  $\mathcal{P}$  of  $D$  is a set of vertex-disjoint paths of  $D$  that cover  $V(D)$ . We say that  $\mathcal{P}$  is an **optimal** path partition if there is no path partition  $\mathcal{P}'$  of  $D$  such that  $|\mathcal{P}'| < |\mathcal{P}|$ . We denote by  $\pi(D)$  the size of an optimal path partition of a digraph  $D$ .

Dilworth [Dilworth 1950] showed that for every transitive acyclic digraph  $D$  we have  $\pi(D) = \alpha(D)$ . Note that this equality is not valid for any digraph; for example, if  $D$  is a directed cycle with 5 vertices, then  $\pi(D) = 1$  and  $\alpha(D) = 2$ . However, Gallai and Milgram [Gallai and Milgram 1960] have shown that  $\pi(D) \leq \alpha(D)$  for every digraph  $D$ .

Greene and Kleitman [Greene and Kleitman 1976] proved a generalization of Dilworth's Theorem described next. Let  $k$  be a positive integer. The  **$k$ -norm** of a path partition  $\mathcal{P}$ , denoted by  $|\mathcal{P}|_k$ , is defined as  $|\mathcal{P}|_k = \sum_{P \in \mathcal{P}} \min\{|P|, k\}$ . We say that  $\mathcal{P}$  is a  **$k$ -optimal path partition** if there is no path partition  $\mathcal{P}'$  such that  $|\mathcal{P}'|_k < |\mathcal{P}|_k$ . We denote by  $\pi_k(D)$  the  $k$ -norm of a  $k$ -optimal path partition of  $D$ . A  **$k$ -partial coloring**  $\mathcal{C}^k$  is a set of  $k$  disjoint stable sets called **color classes** (empty color classes are allowed). The **weight** of a  $k$ -partial coloring  $\mathcal{C}^k$ , denoted by  $\|\mathcal{C}^k\|$ , is defined as  $\|\mathcal{C}^k\| = \sum_{C \in \mathcal{C}^k} |C|$ . We say that  $\mathcal{C}^k$  is an **optimal  $k$ -partial coloring** if there is no  $k$ -partial coloring  $\mathcal{B}^k$  such that  $\|\mathcal{B}^k\| > \|\mathcal{C}^k\|$ . We denote by  $\alpha_k(D)$  the weight of an optimal  $k$ -partial coloring of  $D$ . Given these definitions, what Greene and Kleitman [Greene and Kleitman 1976] showed was that for every transitive acyclic digraph  $D$ , we have  $\pi_k(D) = \alpha_k(D)$ . Note that  $\pi(D) = \pi_1(D)$  and  $\alpha(D) = \alpha_1(D)$ . Thus, Dilworth's Theorem is a particular case of Greene-Kleitman's Theorem in which  $k = 1$ .

As Gallai-Milgram's Theorem extends Dilworth's Theorem, it is a natural question whether Greene-Kleitman's Theorem can be extended to digraphs in general. More precisely, is it true that for every digraph  $D$  we have that  $\pi_k(D) \leq \alpha_k(D)$ ? Linial [Linial 1981] conjectured that the answer for this question is positive.

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<sup>1</sup>Usually  $|P|$  denotes the length of a path (number of arcs), but here it denotes the number of vertices.

**Linial's Conjecture [Linial 1981].** *Let  $D$  be a digraph and  $k$  be a positive integer. Then,  $\pi_k(D) \leq \alpha_k(D)$ .*

Linial's Conjecture remains open, but we know it holds for acyclic digraphs [Saks 1979], bipartite digraphs [Berge 1982], digraphs which contain a Hamiltonian path [Berge 1982],  $k = 1$  [Linial 1978],  $k = 2$  [Berger and Hartman 2008] and  $k \geq \lambda(D) - 3$  [Herskovics 2013]. In this paper we give partial results on Linial's Conjecture for split digraphs.

## 2. Split digraphs

Let  $D$  be a digraph and let  $X \subseteq V(D)$ . We denote by  $D[X]$  the subdigraph of  $D$  induced by  $X$ . A digraph  $D$  is a **split digraph** if there is a partition of  $V(D)$  into two sets  $X$  and  $Y$ , such that  $D[X]$  is a tournament and  $D[Y]$  is a stable set. We shall use the notation  $D[X, Y]$  to indicate that  $D$  is a split digraph with such partition  $\{X, Y\}$ .

In this section we shall prove an approximation to Linial's Conjecture for split digraphs, i. e., that  $\pi_k(D) \leq \alpha_k(D) + 1$  for every split digraph  $D$ , as stated in Theorem 1. For that, we need Rédei's Theorem and Lemmas 1 and 2 below.

**Rédei's Theorem [Rédei 1934].** *Every tournament contains a Hamiltonian path.*

**Lemma 1.** *Let  $D[X, Y]$  be a split digraph. Then,  $\pi_k(D) \leq |Y| + \min\{|X|, k\}$ .*

*Proof.* By Rédei's Theorem, the tournament  $D[X]$  contains a path  $P$  such that  $V(P) = X$ . Let  $\mathcal{P} = \{P\} \cup \{(y) : y \in Y\}$ . Clearly,  $\mathcal{P}$  is a path partition of  $D$  for which  $|\mathcal{P}|_k = \min\{|X|, k\} + |Y|$ . Therefore,  $\pi_k(D) \leq |\mathcal{P}|_k = \min\{|X|, k\} + |Y|$ . ■

**Lemma 2.** *Let  $D[X, Y]$  be a split digraph. Then,  $\alpha_k(D) \geq |Y| + \min\{|X|, k - 1\}$ . Moreover, when  $|X| < k$ , we have that  $\alpha_k(D) = |V(D)|$ .*

*Proof.* First, suppose that  $|X| \leq k - 1$ . Let  $\mathcal{C}^k = \{Y\} \cup \{\{x\} : x \in X\}$ . Note that  $\mathcal{C}^k$  is a  $k$ -partial coloring of  $D$  with  $\|\mathcal{C}^k\| = |V(D)|$ . Therefore,  $\alpha_k(D) = \|\mathcal{C}^k\| = |Y| + |X| = |Y| + \min\{|X|, k - 1\}$  and the result follows. We may assume that  $|X| \geq k$ . Let  $\mathcal{C}^k = \{Y\} \cup \{\{x\} : x \in S\}$ , where  $S \subseteq X$  such that  $|S| = k - 1$ . Clearly,  $\mathcal{C}^k$  is a  $k$ -partial coloring for which  $\|\mathcal{C}^k\| = |Y| + k - 1$ . Therefore,  $\alpha_k(D) \geq \|\mathcal{C}^k\| = |Y| + k - 1 = |Y| + \min\{|X|, k - 1\}$ . ■

**Theorem 1.** *Let  $D[X, Y]$  be a split digraph. Then,  $\pi_k(D) \leq \alpha_k(D) + 1$ .*

*Proof.* The result follows immediately from Lemmas 1 and 2. ■

In Section 2.1 we introduce  $k$ -loose digraphs and show that Linial's Conjecture holds for them and in Section 2.2 we show that it holds for spider digraphs [Hoàng 1985].

### 2.1. $k$ -loose digraphs

A split digraph  $D[X, Y]$  is  **$k$ -loose** if either  $|X| < k$  or there is a  $S \subseteq X$  such that  $|S| = k$  and no vertex  $y \in Y$  is adjacent to every vertex in  $S$ . A split digraph  $D[X, Y]$  that is not  $k$ -loose is called  **$k$ -tight**. We show in this section that Linial's Conjecture holds for every  $k$ -loose digraph (Theorem 2) and for split digraphs such that  $|X| \leq k$  (Theorem 3). For that, we need Lemmas 3 and 4 below.

**Lemma 3.** *Let  $D[X, Y]$  be a split digraph. Then,  $D$  is  $k$ -loose if and only if  $\alpha_k(D) \geq |Y| + \min\{|X|, k\}$ .*

*Proof.* Consider that  $D$  is  $k$ -loose. If  $|X| < k$ , then by Lemma 2,  $\alpha_k(D) = |V(D)| = |Y| + |X| \geq |Y| + \min\{|X|, k\}$ . We may thus assume that  $|X| \geq k$  and there is  $S \subseteq X$  such that  $|S| = k$  and no vertex  $y \in Y$  is adjacent to every vertex in  $S$ . Assume  $S = \{x_1, x_2, \dots, x_k\}$  and let  $\mathcal{C}_0^k = \{C_1, C_2, \dots, C_k\}$  be a  $k$ -partial coloring where  $C_i = \{x_i\}$  for  $i = 1, 2, \dots, k$ . For each  $y \in Y$  choose some vertex  $x_i$  not adjacent to  $y$  (which exists by definition) and add  $y$  in color class  $C_i$ . The  $k$ -partial coloring  $\mathcal{C}^k$  thus obtained has weight  $|Y| + k = |Y| + \min\{|X|, k\}$  as expected.

Conversely, consider that  $\alpha_k(D) \geq |Y| + \min\{|X|, k\}$ . If  $|X| < k$ , then  $D$  is  $k$ -loose by definition. So, we may assume that  $|X| \geq k$  and, whence,  $\alpha_k(D) \geq |Y| + k$ . We conclude that  $\mathcal{C}^k$  must have exactly  $k$  vertices of  $X$ , besides all  $|Y|$  vertices from  $Y$ . Let  $S = \{x : x \in C_i \cap X \text{ for } i = 1, 2, \dots, k\}$ . Since all vertices of  $Y$  belong to  $\mathcal{C}^k$ , then there is no vertex in  $Y$  which is adjacent to every vertex of  $S$ . Therefore,  $D$  is  $k$ -loose. ■

**Theorem 2.** *Let  $D[X, Y]$  be a  $k$ -loose split digraph. Then,  $\pi_k(D) \leq \alpha_k(D)$ .*

*Proof.* By Lemma 3,  $\alpha_k(D) \geq |Y| + \min\{|X|, k\}$ . On the other hand, by Lemma 1  $\pi_k(D) \leq |Y| + \min\{|X|, k\}$  and the result follows. ■

**Lemma 4.** *Let  $D[X, Y]$  be a split digraph such that  $\lambda(D) > |X|$ . Then,  $\pi_k(D) \leq \alpha_k(D)$ .*

*Proof.* If  $\alpha_k(D) = |V(D)|$ , then the result follows trivially. Thus, we may assume that  $\alpha_k(D) < |V(D)|$ . By Lemma 2 we have that  $|X| \geq k$  and also that  $\alpha_k(D) \geq |Y| + \min\{|X|, k - 1\} = |Y| + k - 1$ . Since  $\lambda(D) > |X|$ , there exists a path  $P$  in  $D$  such that  $|P| = |X| + 1$ . Let  $\mathcal{P} = \{P\} \cup \{(v) : v \notin V(P)\}$ . Clearly,  $\mathcal{P}$  is a path partition of  $D$  and  $|\mathcal{P}|_k = |Y| + k - 1$ . Therefore,  $\pi_k(D) \leq |\mathcal{P}|_k \leq \alpha_k(D)$ . ■

**Theorem 3.** *Let  $D[X, Y]$  be a split digraph such that  $|X| \leq k$ . Then,  $\pi_k(D) \leq \alpha_k(D)$ .*

*Proof.* If  $D$  is  $k$ -loose, then the result follows by Theorem 2. So, we may assume that  $D$  is not  $k$ -loose. Hence,  $|X| = k$  and there exists a vertex  $y \in Y$  which is adjacent to every vertex of  $X$ . Therefore,  $D[X \cup \{y\}]$  is a tournament and by Rédei's Theorem it has a Hamiltonian path  $P$  such that  $|P| = |X| + 1$ . As  $P$  is a path in  $D$  as well, we conclude that  $\lambda(D) \geq |X| + 1$  and the result follows by Lemma 4. ■

## 2.2. Spider digraphs

We denote by  $\mathcal{N}(v)$  the set of vertices that are adjacent to  $v \in V(D)$  (regardless the direction of the arcs). A split digraph  $D[X, Y]$  is **spider** [Hoàng 1985] if (i)  $|X| = |Y| \geq 2$ ; and (ii) there exists a bijective function  $f : X \rightarrow Y$  such that either  $\mathcal{N}(x) = \{f(x)\}$  for all  $x \in X$  (in this case, we say that  $D$  is a **thin** spider) or  $\mathcal{N}(x) = Y - f(x)$  for all  $x \in X$  (in this case, we say that  $D$  is a **thick** spider). Note that thin spider digraphs are  $k$ -loose, but thick spider digraphs are  $k$ -tight, as long as  $|X| > k$ . The following theorem shows that Linial's Conjecture holds for spider digraphs.

**Theorem 4.** *Let  $D[X, Y]$  be a spider digraph. Then,  $\pi_k(D) \leq \alpha_k(D)$ .*

*Proof.* Let  $\ell = |X| = |Y|$ . If  $\ell \leq k$ , then the result follows by Theorem 3. Thus, we may assume that  $|X| > k$ . Clearly,  $\pi_k(D) \leq |V(D)|$  and we deduce that  $\alpha_k < |V(D)|$ . If  $D$  is a thin spider digraph, whence  $k$ -loose, the result follows by Theorem 2. Therefore, we may assume that  $D$  is a thick spider graph. Since  $D[X]$  is a tournament, by Rédei's Theorem, there exists a path  $P$  such that  $V(P) = X$ . Let  $P = (x_1, x_2, \dots, x_\ell)$ . Since  $D$  is a thick spider digraph, there exists one single vertex  $y_i \in Y$  that is not adjacent to  $x_i$ , for  $i = 1, \dots, \ell$ . Note that if  $\lambda(D) > |X|$ , then the result follows by Lemma 4. So we may assume that  $\lambda(D) \leq |X|$ .

Let  $Px_i$  denote the subpath  $(x_1, x_2, \dots, x_i)$  and let  $x_iP$  denote the subpath  $(x_i, x_{i+1}, \dots, x_\ell)$ . We denote by  $W \circ Q$  the concatenation of two paths  $W$  and  $Q$ .

**Claim 1:** *If  $x_i \in X$ ,  $y_j \in Y$  and  $i < j$ , then  $(x_i, y_j) \in A(D)$ .*

We prove this claim by induction on  $i$ . If  $i = 1$ , assume by contradiction that  $(y_j, x_1) \in A(D)$ ; then  $P' = (y_j, x_1) \circ P$  is a path in  $D$  such that  $|P'| = |X| + 1$ , a contradiction. Hence,  $(x_1, y_j) \in A(D)$ . Consider now  $i > 1$ . Recall that  $y_j$  is adjacent to every vertex in  $X - \{x_j\}$ . Thus,  $y_j$  is adjacent to every vertex of  $V(Px_i)$ . By induction hypothesis, we have  $(x_{i-1}, y_j) \in A(D)$ . Suppose by contradiction that  $(y_j, x_i) \in A(D)$ . Then, there is a path  $P' = Px_{i-1} \circ (x_{i-1}, y_j, x_i) \circ x_iP$  such that  $|P'| = |X| + 1$ , a contradiction. Therefore,  $(x_i, y_j) \in A(D)$ . This completes the proof of Claim 1.

**Claim 2:** *If  $x_i \in X$ ,  $y_j \in Y$  and  $j < i$ , then  $(y_j, x_i) \in A(D)$ .*

We omit the proof of Claim 2, as it is analogous to that of Claim 1.

We claim that both  $P_0 = (x_1, y_2, x_3, y_4, \dots)$  and  $P_1 = (y_1, x_2, y_3, x_4, \dots)$  are paths in  $D$ . By Claim 1 we have that  $(x_i, y_{i+1}) \in A(D)$  for  $i = 1, 3, \dots$ , and by Claim 2 we have that  $(y_j, x_{j+1}) \in A(D)$  for  $j = 2, 4, \dots$ . Hence  $P_0$  is a path in  $D$ . The proof is analogous for  $P_1$ . Clearly,  $\mathcal{P} = \{P_0, P_1\}$  is a path partition of  $D$ . Moreover,  $|P_0| = |P_1| = \ell$  and  $|\mathcal{P}|_k = 2 \min\{\ell, k\} = 2k$ . Since  $|X| > k$ , we have that  $\min\{\ell, k\} = k \leq |X| - 1 = |Y| - 1$ . Thus,  $|\mathcal{P}|_k = 2k \leq k + |Y| - 1$ . On the other hand, by Lemma 2,  $\alpha_k(D) \geq |Y| + \min\{|X|, k - 1\} = |Y| + k - 1$ . Therefore,  $\pi_k(D) \leq |\mathcal{P}|_k = 2k \leq |Y| + k - 1 \leq \alpha_k(D)$ . ■

### 3. Conclusion

We showed that Linial's Conjecture holds for  $k$ -loose digraphs and for some subclasses of  $k$ -tight digraphs, namely those with  $|X| = k$  and the thick spider digraphs. It is easy to see that for  $k$ -tight digraphs,  $\alpha_k(D) = |Y| + k - 1$ . Therefore, it is clear that any approach to prove Linial's Conjecture for  $k$ -tight digraphs must involve finding a path partition with  $k$ -norm less than or equal to  $|Y| + k - 1$ . We are currently working on this idea.

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